

**MAL 609 Basic Computer Science
Tutorial Sheet -I**

(Asymptotic)

1. Suppose $T_1(n) = \Omega(f(n))$ and $T_2(n) = \Omega(g(n))$. Which of the following statements are true?
 - (a) $T_1(n) + T_2(n) = \Omega(\max\{f(n), g(n)\})$.
 - (b) $T_1(n)T_2(n) = \Omega(f(n)g(n))$.
 - (c)
2. If $f(n)$ is a k^{th} degree polynomial, then show that
 - (a) $f(n) = O(n^k)$.
 - (b) $f(n) = \Omega(n^k)$.
3. Order the following functions by growth rate: $N, \sqrt{N}, N^{1.5}, N^2, N \log N, N \log \log N, N \log^2 N, N \log(N^2), 2/N, 2^N, 2^{N/2}, 37, N^2 \log N, N^3$. Indicate which functions grow at the same rate.
4. Suppose $T_1(N) = O(f(N))$ and $T_2(N) = O(f(N))$. Which of the following are true ?
 - a. $T_1(N) + T_2(N) = O(f(N))$
 - b. $T_1(N) - T_2(N) = o(f(N))$
 - c. $T_1(N) / T_2(N) = O(1)$
 - d. $T_1(N) = O(T_2(N))$.
5. Which functions grows faster: $N \log N$ or $N^{1+\epsilon/\sqrt{\log N}}, \epsilon > 0$?
6. Prove that for any constant, $k, \log^k N = o(N)$.
7. Find two functions $f(N)$ and $g(N)$ such that neither $f(N) = O(g(N))$ nor $g(N) = O(f(N))$.
8. For each function $f(n)$ and time t in the following table, determine the largest size n of a problem that can be solved in time t assuming that the algorithm to solve the problem takes $f(n)$ microseconds.

	1 second	1 hour	1 Month	1 century
$\log n$				
\sqrt{n}				
N				
$n \log n$				
n^2				
n^3				
2^n				
$n!$				

9. Show that $\log^3 n$ is $o(n^{1/3})$.
10. Show that $\sum_{i=1}^n i^2$ is $O(n^3)$.
11. Consider the following functions of n :

$$f_1(n) = n^2$$

$$f_2(n) = n^2 + 1000n$$

$$f_3(n) = n, \text{ if } n \text{ is odd}$$

$$n^3, \text{ if } n \text{ is even.}$$

$$f_4(n) = n, \text{ if } n \leq 100$$

$$n^3, \text{ if } n > 100$$
 Indicate for each distinct pair i and j whether $f_i(n)$ is $O(f_j(n))$ and whether $f_j(n)$ is $\Omega(f_i(n))$.
12. Consider the following functions of n :

$$g_1(n) = n^2 \text{ for even } n \geq 0$$

$$n^3 \text{ for odd } n \geq 1$$

$$g_2(n) = n, 0 \leq n \leq 100$$

$$n^3, \text{ for } n > 100$$

$$g_3(n) = n^{2.5}$$

Indicate for each distinct pair i and j whether $g_i(n)$ is $O(g_j(n))$ and whether $g_j(n)$ is $\Omega(g_i(n))$.

13. Give, using “bog oh” notation, the worst case running times of the following procedures as a function of n.

- a. **procedure** *matmpy* (n : integer);
 var
 i, j, k : integer;
begin
 for I = 1 to n
 for j = 1 to n do **begin**
 C[I,j] = 0
 for k = 1 to n do
 C[I, j] = C[I,j] + A[I, k]*B[k,j]
 end
 end
end
- b. **procedure** *mystery* (n: integer)
 var
 i,j,k : integer;

begin
 for i = 1 to n-1 do
 for j= i+1 to n do
 for k = 1 to j do
 {some statement requiring O(1) time}
 end
end
- c. **procedure** *veryodd* (n : integer);
 var
 i, j, x, y : integer;
begin
 for i = 1 to n
 if *odd*(i) **then begin**
 for j = i to n do **begin**
 x = x+1;

 for j= 1 to ido
 y = y +1;
 end
 end
end
- d. **function** *recursive*(n:integer) : integer;
begin
 if n <= 1 **then**
 return(1)
 else
 return(*recursive*(n-1) + *recursive*(n-1))
end

14. Show that the following statements are true.

- 17 is $O(1)$.
 - $n(n-1)$ is $O(n^2)$
 - $\sum_{i=1}^n i^k$ is $O(n^{k+1})$ and $\Omega(n^{k+1})$ for integer k.
 - If $P(x)$ is any k^{th} degree polynomial with a positive leading coefficient, then $P(n)$ is $O(n^k)$ and $\Omega(n^k)$.
15. Order the following functions by growth rate: n , \sqrt{n} , $\log n$, $\log \log n$, $\log^2 n$, $n/\log n$, $\sqrt{n \log^2 n}$, $(1/3)^n$, $(3/2)^n$, 17.