

# Zero Sum Games — How to find sol<sup>n</sup> to them using L.P.

Eg. Game — Prisoner's Dilemma.

Police → A, B  
(who has committed the crime)

| A \ B |                              | Silent     | Other person has committed crime |
|-------|------------------------------|------------|----------------------------------|
|       |                              | Silent     | Other person committed crime     |
| A     | Silent                       | $(-1, -1)$ | $(3, 0)$                         |
|       | Other person committed crime | $(0, -3)$  | $(-2, -2)$                       |

Equilibrium strategy

Ques: What will be best action for (A, B) = (Silent, Silent)  
Penalty = (-1, -1) ← Not an Equilibrium

• If A betrays = { A - Other person ---, B - Silent }  
Penalty = (0, -3)

• If B betrays = { , }  
Penalty = (-3, 0)

\* Equilibrium (Nash Equilibrium) = (Other person, Other person)  
(-2, -2)

Def<sup>n</sup> Equilibrium A strategy pair  $(x, y)$  for players  $(A, B)$  is Nash Equilibrium if

(i) for  $x' \neq x$  gain of A at  $(x', y) \not\geq$  gain of A at  $(x, y)$

(A is not better-off by choosing strategy  $x'$  instead of  $x$ )

(ii) for  $y' \neq y$  gain of B at  $(x, y') \not\geq$  gain of B at  $(x, y)$

Eg. Rock - Paper - Scissor game

|                                     | A \ B          | R <sub>0</sub> | P <sub>e</sub> | S <sub>c</sub> |
|-------------------------------------|----------------|----------------|----------------|----------------|
| R <sub>0</sub> $\gg$ S <sub>c</sub> | R <sub>0</sub> | 0              | -1             | 1              |
| S <sub>c</sub> $>$ P <sub>a</sub>   | P <sub>a</sub> | 1              | 0              | -1             |
| P <sub>a</sub> $>$ R <sub>0</sub>   | S <sub>c</sub> | -1             | 1              | 0              |

What A wins

|                | A \ B          | R <sub>0</sub> | P <sub>e</sub> | S <sub>c</sub> |
|----------------|----------------|----------------|----------------|----------------|
| R <sub>0</sub> | R <sub>0</sub> | 0              | 1              | -1             |
| P <sub>a</sub> | P <sub>a</sub> | -1             | 0              | 1              |
| S <sub>c</sub> | S <sub>c</sub> | 1              | -1             | 0              |

What B wins

- Each row / column has  $(+1, -1, 0)$

- (Whatever A wins) + (Whatever B wins) = 0

CLAIM: There is no (Pure / Det) Nash Eq in above game.

Proof: Try all possibilities

An example of randomized strategy (mixed strategy).

$$\left( \begin{array}{c|ccc} A & R_0 & P_a & S_c \\ \hline \text{w.p.} & 0.2 & 0.3 & 0.5 \end{array} \right)$$

$$\left( \begin{array}{c|ccc} B & R_0 & P_a & S_c \\ \hline \text{w.p.} & 1/2 & 1/2 & 0 \end{array} \right)$$

|       |                | B              |                |                |
|-------|----------------|----------------|----------------|----------------|
|       |                | R <sub>0</sub> | P <sub>a</sub> | S <sub>c</sub> |
| 0.2 — | R <sub>0</sub> | 0              | -1             | 1              |
| 0.3 — | P <sub>a</sub> | 1              | 0              | -1             |
| 0.5 — | S <sub>c</sub> | -1             | 1              | 0              |

What A wins

Expected  
Payoff  
of  
A

$$\frac{0.2}{2} (0) + \frac{(0.3)}{2} (1) + \dots$$

Zero Sum Games : Games in which for any choice of strategies for players (A, B)

$$(\text{Payoff of A}) + (\text{Pay-off of B}) = 0$$

↑ To express them, only one matrix suffices

$$(P) = (p_{ij}) \quad \text{— pay off of A}$$

$$-P \quad \text{— pay off of B}$$

Mixed strategy of A :  $x = (x_1 \dots x_m)$  with  $\sum_{i=1}^m x_i = 1, x \geq 0$

Mixed strategy of B :  $y = (y_1 \dots y_n)$  with  $\sum_{j=1}^n y_j = 1, y \geq 0$

Expected Pay-off of A :

$$= \sum_{i,j} p_{ij} \cdot \text{Prob}[\text{Alice plays } i] \cdot \text{Prob}[\text{Bob plays } j]$$

$$= \sum_{i,j} p_{ij} \cdot x_i \cdot y_j = x^T P y$$

Expected Pay-off of B :

$$= \sum_{i,j} (-p_{ij}) \cdot \text{Prob}[\text{Alice plays } i] \cdot \text{Prob}[\text{Bob plays } j]$$

$$= \sum_{i,j} (-p_{ij}) \cdot x_i \cdot y_j = (-x^T P y)$$

Ques : How to define Equilibrium when we have mixed/  
randomized strategies.

Def<sup>n</sup> A mixed strategy pair  $(\tilde{x}, \tilde{y})$  for players (A,B) is  
Nash Equilibrium if

$$(i) \text{ for } x' \neq \tilde{x} \quad \begin{array}{c} \text{Exp} \\ \text{Payoff of A at} \\ (x', \tilde{y}) \end{array} \neq \begin{array}{c} \text{Exp} \\ \text{Payoff of A} \\ \text{at } (\tilde{x}, \tilde{y}) \end{array}$$

$$(ii) \text{ for } y' \neq \tilde{y} \quad \begin{array}{c} \text{Exp} \\ \text{Payoff of B at} \\ (\tilde{x}, y') \end{array} \neq \begin{array}{c} \text{Exp} \\ \text{Payoff of B} \\ \text{at } (\tilde{x}, \tilde{y}) \end{array}$$

H.W. 1 Suppose A has declared its mixed strategy

$$x = (x_1, \dots, x_m) \quad \left( \sum x_i = 1, x \geq 0 \right).$$

Then, how can we find best mixed strategy for y.

Hint: Very simple L.P.

x is known / declared

— Find best "y" for B

H.W. 2 Compute Dual of LP you get in H.W. 1