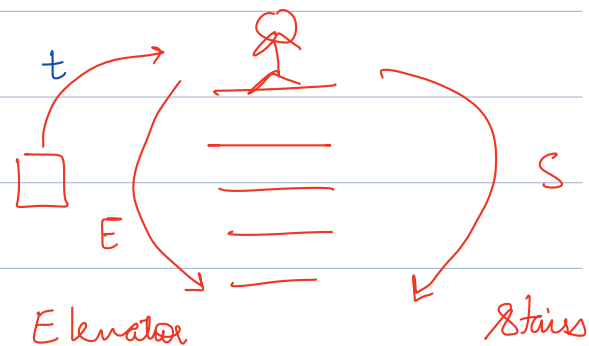


Online Algo - Elevator Puzzle

optimize :

Time to reach Ground floor



Best Offline Algo (when t was known)

If $t + E \leq S$: Use elevator
else : use stairs

} time = $\text{MIN}(t + E, S)$

t = Time taken by elevator to reach top floor (Unknown)
 S = Time by steps to reach ground floor.

Online
Setting

Unknown : Time " t " lift will take to reach Top floor

Best online strategy = ?

Natural Algo

Wait for " α " seconds,
if lift arrives, then take
it, else take Stairs

$$\text{Time} = \begin{cases} t + E & t \leq \alpha \\ \alpha + S & t > \alpha \end{cases}$$

Algo 1 $\alpha = S$

Wait for "S" seconds.

If elevator comes : Take elevator
else : Take stairs

Best Offline Algo.

$$\text{Time} = \begin{cases} \underline{t + E} & t \leq S \\ S + S & t > S \end{cases}$$

$$\min(\underline{t + E}, \underline{S})$$

$C_2(\text{Algo 1}) = \max_t \left(\frac{\text{Time taken by Algo 1}}{\min(t + E, S)} \right)$ — How far are we from best offline algo.

Scenario 1 $t \leq S - E$

$$\text{ratio} = \frac{t + E}{t + E} = 1$$

Scenario 2 $t \in [S - E, S]$

$$\text{ratio} = \frac{t + E}{S}$$

Scenario 3 $t \geq S$

$$\text{ratio} = \frac{2S}{S}$$

$$C_2 = 2$$

Algo 2 $\alpha = S - E$

Best Wait for "S-E" seconds.

If elevator comes : Take elevator
else : Take stairs

H.W(1)

$$C_2 = \left(2 - \frac{E}{S} \right)$$

$$t \leq S - \epsilon$$

$$t \geq S - \epsilon$$

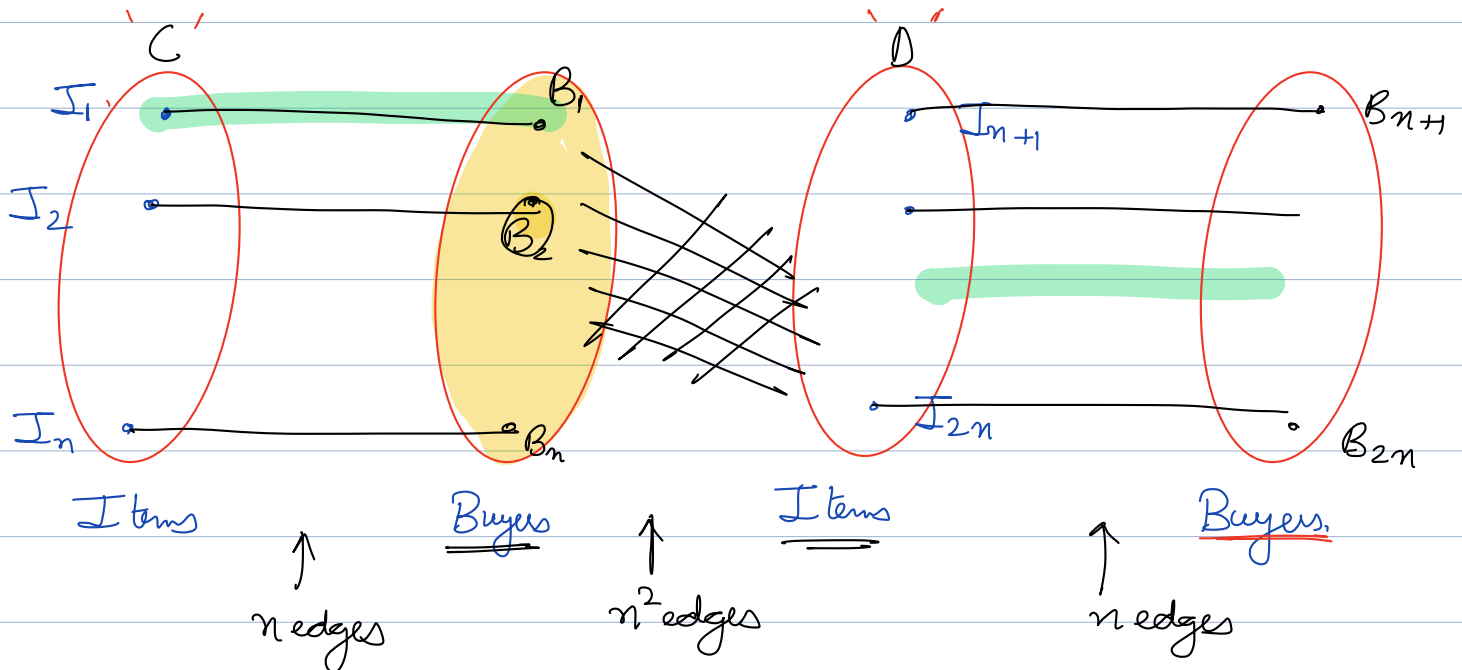
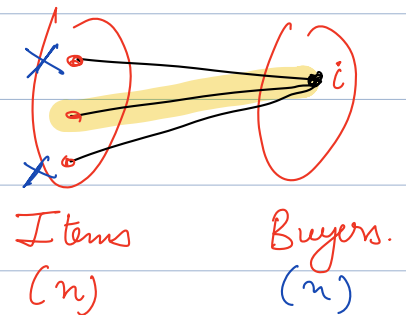
H.W. 3 For any det algo for Elevator problem $C \geq 2 - \frac{\epsilon}{5}$

Online Matching

Algo

For $i = 1$ to n :

match buyer B_i "randomly" to any of available goods
he/she wants to buy.



Offline scenario - what is $|\text{maximum-matching}| = 2n$

Online Scenario :

ques - what is expected size of matching in online setting.

Random variable $X_i = \begin{cases} 1 & \text{if buyer } B_i \text{ is matched in set } C \\ 0 & \text{o/w} \end{cases}$
For $i \leq n$

CLAIM : $\left(\sum_{i=1}^n X_i + n \right)$ - size of online matching.

$$\text{Exp size of online matching} = \sum_{i=1}^n E(X_i) + n$$

$$E(X_1) = \text{Prob}(B_1 \text{ is matched in set } C) = \frac{1}{n+1}$$

$$E(X_2) = \text{Prob}(B_2 \text{ is matched in set } C) = \begin{cases} \frac{1}{n+1} & \text{if first buyer matched in set } C \\ \frac{1}{n} & \text{o/w} \end{cases}$$

$$\leq \frac{1}{n}$$

$$E(X_3) \leq \frac{1}{n-1}$$

$$E(X_n) \leq \frac{1}{2}$$

$$\begin{aligned} \text{Exp size of our matching} &\leq \left(\frac{1}{n+1} + \frac{1}{n} + \dots + \frac{1}{2} \right) + n \\ &= O(\log n) + n \end{aligned}$$

$$\underline{\underline{C_r(\text{rand algo})}} \leq \sup_n \left(\frac{n + O(\log n)}{2n} \right) = \frac{1}{2} \quad \text{--- (i)}$$

$$C_r(\text{rand algo}) \geq \frac{1}{2} \quad \text{since we compute maximal matching.} \quad \text{--- (ii)}$$

$$C_r(\text{rand algo}) = \frac{1}{2}$$