

* Quiz + Discussion (2-3 PM)

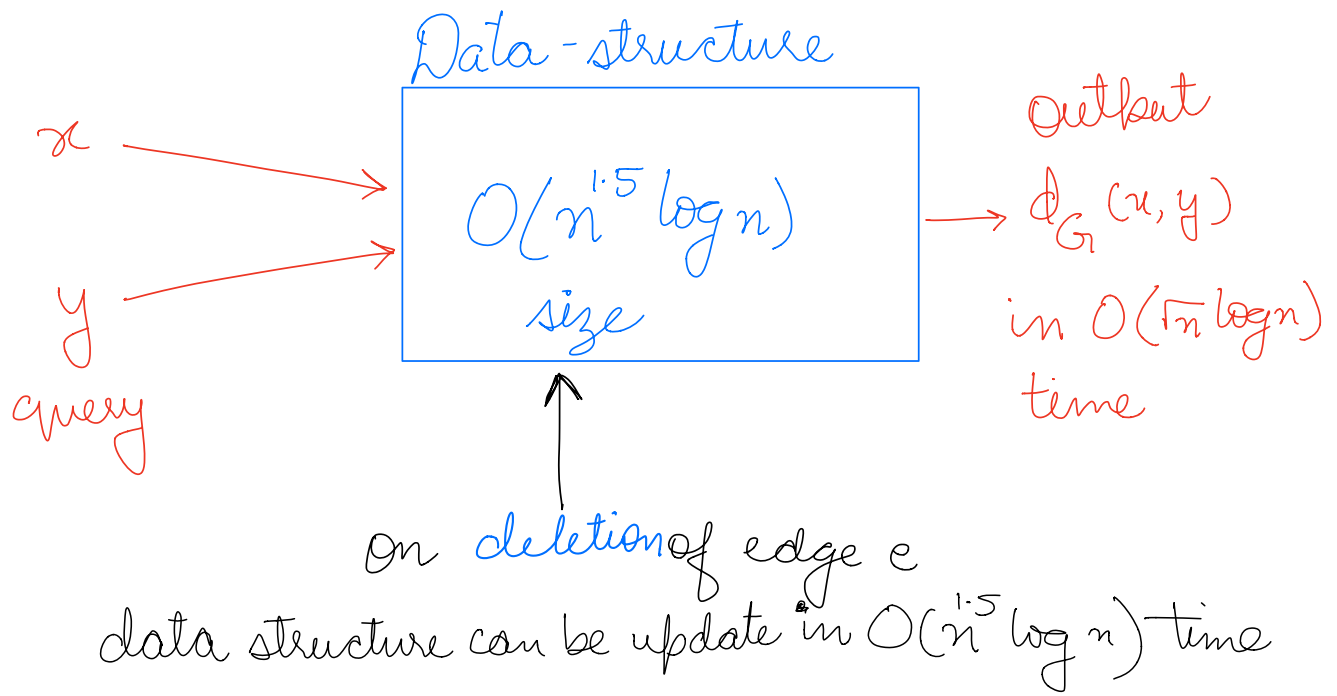
* Student talk - Decremental SSSP

Application of decremental SSSP:

Problem Statement :

Maintain a oracle (data-structure) of $O(n^{1.5} \log n)$ size *decrementally* (i.e. for a graph undergoing edge *deletions*), so that :

- (i) after every edge *deletion*, the oracle can be updated quickly (in amortized $O(n^{1.5} \log n)$ time.)
- (ii) given any query pair (x, y) the oracle reports $d_G(x, y)$ in $O(\sqrt{n} \log n)$ time.



We mimic the algorithm taught in last class:

Data-Structure

- ① For each $v \in V$
we maintain out-BFS $_{\sqrt{n}}(v)$
 - ② $R \leftarrow$ Random set of size " $16\sqrt{n} \log n$ "
For each $v \in R$,
we maintain in-bfs (v) & out-bfs (v)
- (decrementally)

Total update time:

$$\begin{aligned}
 & |V| \cdot O(m\sqrt{n}) + |R| \cdot O(mn) \\
 & = O(mn\sqrt{n} \log n) \Rightarrow \text{Amortized update time} \\
 & \quad \quad \quad = \underline{O(n\sqrt{n} \log n)}
 \end{aligned}$$

How to find $d_G(x, y)$ for query pair $(x, y) \in V \times V$?

① If $y \in \text{OUT-BFS}_{\sqrt{n}}(x)$	② If $y \notin \text{OUT-BFS}_{\sqrt{n}}(x)$
Report $d_G(x, y)$ from tree of x in $O(1)$ time.	$\Rightarrow d_G(x, y) > \sqrt{n}$  Report $\min_{r \in R} \{d_G(x, r) + d_G(r, y)\}$

worst case query time

$$= O(|R|) = O(\sqrt{n} \log n)$$

Question: How to incrementally / decrementally maintain a $+2$ additive spanner?

Recall. (lecture 2)

Algorithm to compute "+2" additive spanner

- ① Initialize $H = (V, E_H = \emptyset)$
- ② $R \leftarrow$ uniform random subset of V of size $O(\sqrt{n})$.
- ③ For each $w \in R$:
Add edges of $\text{BFS}(w, G)$ to E_H .
- ④ For each $v \in V(G)$ with $\deg(v, G) \leq 4\sqrt{n} \log n$:
Add to E_H all edges incident to v in G .
- ⑤ Return H

In order to dynamically maintain +2 additive spanner, we just need to

- ① store in H edges incident to all vertices of current degree $\leq \sqrt{n}$.
- ② dynamically maintain $\text{BFS}(z)$, for $z \in R$

Total time will be $= O(m \cdot n \cdot |R|)$

$$= O(m \cdot n \cdot \sqrt{n} \log n)$$

in incremental/decremental setting.