

Question

ls more; more or less

Consider the DFA M

$$= \{Q, \Sigma, \{q_1\}, \delta, F\}$$

where all symbols have their usual meanings.

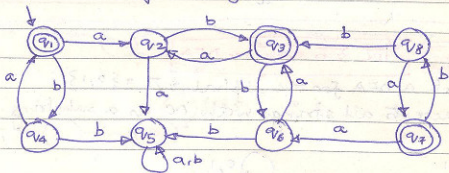
$$Q = \{q_1, q_2, q_3, q_4, q_5, q_6, q_7, q_8\}$$

$$F = \{q_1, q_3, q_7\} \text{ and } \delta \text{ is}$$

given by the table to the right. Using this info, construct an equivalent minimum-state DFA.

State	a	b
q_1	q_2	q_4
q_2	q_5	q_3
q_3	q_2	q_6
q_4	q_1	q_5
q_5	q_5	q_5
q_6	q_3	q_5
q_7	q_6	q_8
q_8	q_7	q_3

'Drawn nicely', the given DFA looks like:



Following the k-equivalence method:

$$P_0 = \{ \underbrace{\{q_1, q_3, q_7\}}_{C_0}, \underbrace{\{q_2, q_4, q_5, q_6, q_8\}}_{C_1} \}$$

		C_0	C_1	C_1	C_1	C_1	C_1	C_0
$x=a$		C_1	C_1	C_1	C_1	C_1	C_1	C_0
$x=b$		C_1	C_1	C_1	C_0	C_1	C_1	C_0
		C_0	C_2	C_3	C_4	C_4	C_5	C_5

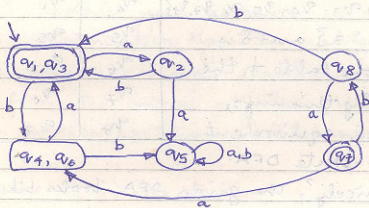
$$P_1 = \{ \underbrace{\{q_1, q_3, q_7\}}_{C_0}, \underbrace{\{q_2\}}_{C_2}, \underbrace{\{q_4, q_6\}}_{C_3}, \underbrace{\{q_5\}}_{C_4}, \underbrace{\{q_8\}}_{C_5} \}$$

		C_0	C_2	C_3	C_4	C_5
$x=a$		C_2	C_2	C_3	C_0	C_0
$x=b$		C_3	C_3	C_5	C_4	C_4
		C_6	C_7	C_3		

$$P_2 = \left\{ \underbrace{\{q_1, q_5\}}_{C_6}, \underbrace{\{q_7\}}_{C_7}, \underbrace{\{q_2\}}_{C_2}, \underbrace{\{q_4, q_6\}}_{C_3}, \underbrace{\{q_5\}}_{C_4}, \underbrace{\{q_8\}}_{C_5} \right\}$$

$$\begin{array}{l}
 x \geq a \quad C_2 \quad C_2 \\
 x = b \quad C_3 \quad C_3 \\
 \underbrace{\hspace{2cm}}_{C_6} \qquad \qquad \underbrace{\hspace{2cm}}_{C_3}
 \end{array}$$

$P_3 = P_2 \implies \text{stop!}$



ALGORITHM FSM_matcher

$q := 0$; compute_ δ ();

FOR $i := 1$ TO n DO /* $\mathcal{O}(n)$ */

$q := \delta(q, T[i])$;

 IF $q == m$ THEN print shift= $i - m$;

ALGORITHM compute_ δ ()

FOR $q := 0$ TO m DO /* $\mathcal{O}(m)$ */

FOREACH $a \in \Sigma$ DO /* $\mathcal{O}(|\Sigma|)$ */

$k := \min(m + 1, q + 2)$; /*no jump-off*/

 REPEAT $k := k - 1$; /* $\leq \mathcal{O}(m)$ times*/

 UNTIL $P[1..k]$ is a suffix of $P[1..q]a$

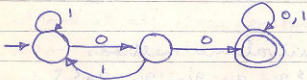
 /* $\leq m$ character comparisons*/

- FSM: $\mathcal{O}(n + m^3 |\Sigma|)$

Question

निम्न स्तरीय $\sqrt{4.4.1} \dots$ DFA

construct a DFA for the alphabet $\Sigma = \{0, 1\}$
which accepts all strings with 00 as a substring.



Question

Even the Introvert must engage
in Regular Expression

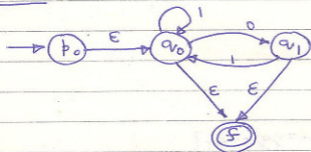
Consider a language over $\{0, 1\}$: all strings with
00 as a substring. Construct a regular
expression for the complement of this
language.

It is easy to go via a DFA, and the state minimisation method.

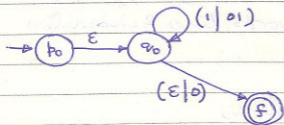
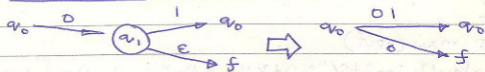


Step 1: Eliminate trap states.

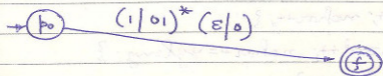
Step 2: New start state, new final state



Eliminate q_1 :



Eliminate q_0



Question

one for the Smart A-lex.../*No Comments*/

What precisely does the following lex code print out, given a few lines of text input, terminated by a Control-D? There will be no marks for an imprecise answer. It is a digital world, you know!

```
%{
```

```
int i, j, k;
```

```
%{
```

```
%%
```

```
\n {k++; i++;}
```

```
[^ \n\t] {j++; i+=yyleng;}
```

```
• {i++;}
```

```
%%
```

```
int main(void)
```

```
{ yylex(); printf("%d\t%d\t%d\n", j, k, i); return 0; }
```

This counts the no of words, lines & characters.

Better written code: -

```
%{
```

```
int nchar, nword, nline;
```

```
%{
```

```
%%
```

```
\n {nline nline++; nchar++;}
```

```
[^ \n\t] {nword++; nchar+=yyleng;}
```

```
• {nchar++;}
```

```
%%
```

```
int main(void)
```

```
{ yylex(); printf("%d\t%d\t%d\n", nchar, nword, nline); return 0; }
```