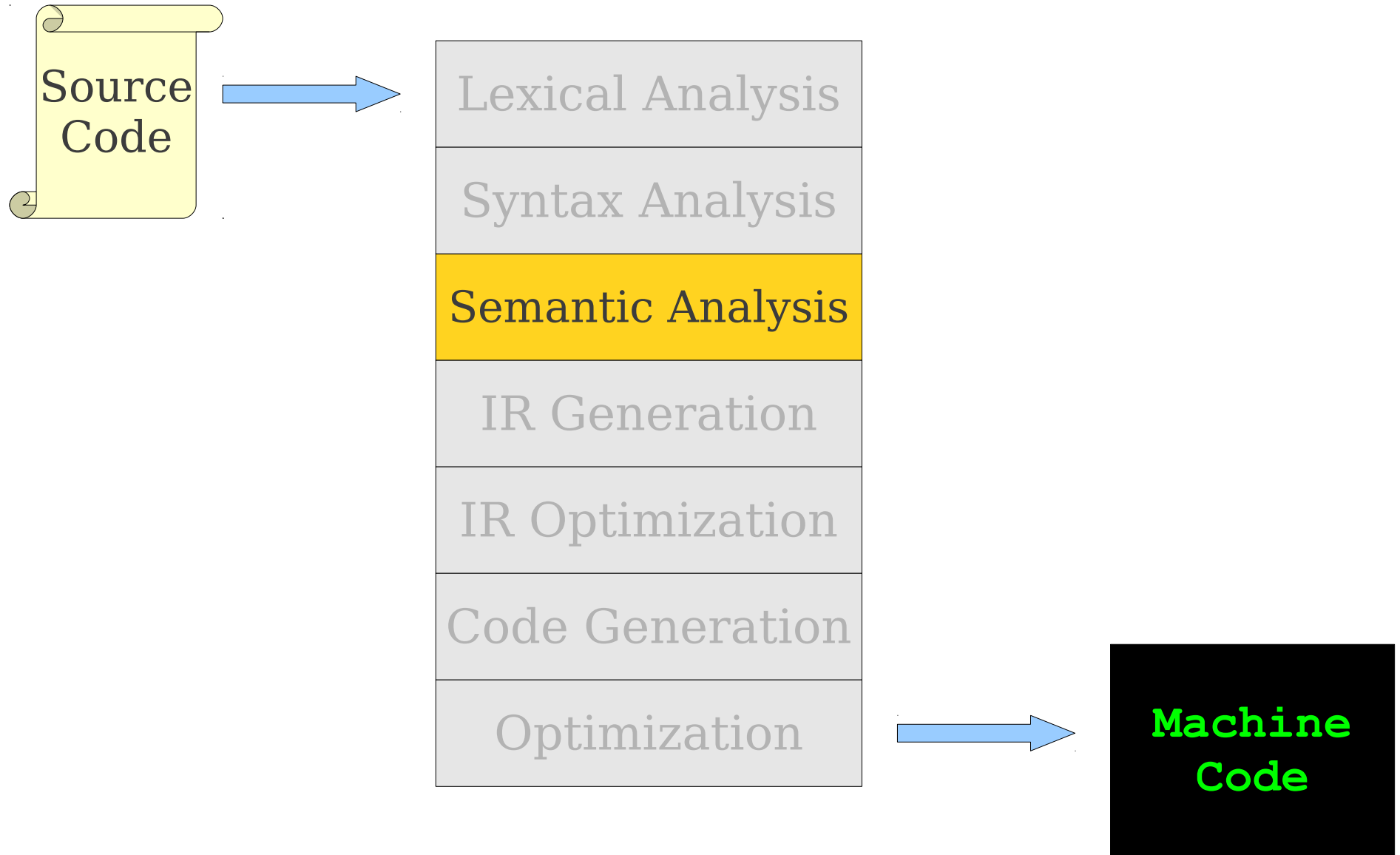


# Type-Checking

# Where We Are



# Review from Last Time

```
class MyClass implements MyInterface {
    string myInteger;

    void doSomething() {
        int[] x;
        x = new string;

        x[5] = myInteger * y;
    }
    void doSomething() {

    }
    int fibonacci(int n) {
        return doSomething() + fibonacci(n - 1);
    }
}
```

# Review from Last Time

```
class MyClass implements  
    string myInteger;
```

**MyInterface** {

Interface not  
declared

```
void doSomething() {  
    int[] x;
```

```
    x = new string;
```

Can't multiply  
strings

Wrong type

```
    x[5] => myInteger * y;
```

Variable not  
declared

```
void doSomething() {
```

Can't redefine  
functions

```
}  
int fibonacci(int n) {  
    return doSomething() + fibonacci(n - 1);  
}
```

Can't add void

```
}
```

No main function

# Review from Last Time

```
class MyClass implements MyInterface {  
    string myInteger;
```

```
    void doSomething() {  
        int[] x;
```

```
        x = new string;
```

```
        x[5] → myInteger * y;
```

```
    }  
    void doSomething() {
```

```
    }  
    int fibonacci(int n) {  
        return doSomething() + fibonacci(n - 1);  
    }
```

```
}
```

Wrong type

Can't multiply strings

Variable not declared

Can't redefine functions

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No main function

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class MyClass implements MyInterface {  
    string myInteger;
```

```
    void doSomething() {  
        int[] x;
```

```
        x = new string;
```

```
        x[5] → myInteger * y;
```

```
    }
```

```
void doSomething() {
```

```
}
```

```
int fibonacci(int n) {
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```
    return doSomething() + fibonacci(n - 1);
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}
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```
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        int[] x;
```

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        x = new string;
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```
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```
    }
```

```
    void doSomething() {
```

```
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int fibonacci(int n) {
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```
    return doSomething() + fibonacci(n - 1);
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}
```

Can't multiply strings

Wrong type

Can't add void



# What Remains to Check?

- **Type errors.**
- Today:
  - What are types?
  - What is type-checking?
  - A type system for Decaf.

# What is a Type?

- This is the subject of some debate.
- To quote Alex Aiken:
  - **“The notion varies from language to language.**
  - The consensus:
    - A set of values.
    - A set of operations on those values”
- **Type errors** arise when operations are performed on values that do not support that operation.

# Types of Type-Checking

- **Static type checking.**
  - Analyze the program during compile-time to prove the absence of type errors.
  - Never let bad things happen at runtime.
- **Dynamic type checking.**
  - Check operations at runtime before performing them.
  - More precise than static type checking, but usually less efficient.
  - (Why?)
- **No type checking.**
  - Throw caution to the wind!

# Type Systems

- The rules governing permissible operations on types forms a **type system**.
- **Strong type systems** are systems that never allow for a type error.
  - Java, Python, JavaScript, LISP, Haskell, etc.
- **Weak type systems** can allow type errors at runtime.
  - C, C++

# Type Wars

- *Endless* debate about what the “right” system is.
- Dynamic type systems make it easier to prototype; static type systems have fewer bugs.
- Strongly-typed languages are more robust, weakly-typed systems are often faster.

# Type Wars

- *Endless* debate about what the “right” system is.
- Dynamic type systems make it easier to prototype; static type systems have fewer bugs.
- Strongly-typed languages are more robust, weakly-typed systems are often faster.
- **I'm staying out of this!**

# Our Focus

- Decaf is typed **statically** and **weakly**:
  - Type-checking occurs at compile-time.
  - Runtime errors like dereferencing `null` or an invalid object are allowed.
- Decaf uses **class-based inheritance**.
- Decaf distinguishes primitive types and classes.

# Typing in Decaf



# Static Typing in Decaf

- Static type checking in Decaf consists of two separate processes:
  - Inferring the type of each expression from the types of its components.
  - Confirming that the types of expressions in certain contexts matches what is expected.
- Logically two steps, but you will probably combine into one pass.

# An Example

```
while (numBitsSet(x + 5) <= 10) {  
  
    if (1.0 + 4.0) {  
        /* ... */  
    }  
  
    while (5 == null) {  
        /* ... */  
    }  
  
}
```

# An Example

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while (numBitsSet(x + 5) <= 10) {  
  
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    }  
}
```

Well-typed  
expression with  
wrong type.

# An Example

```
while (numBitsSet(x + 5) <= 10) {  
  
    if (1.0 + 4.0) {  
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# An Example

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    }  
  
}
```

Expression with  
type error



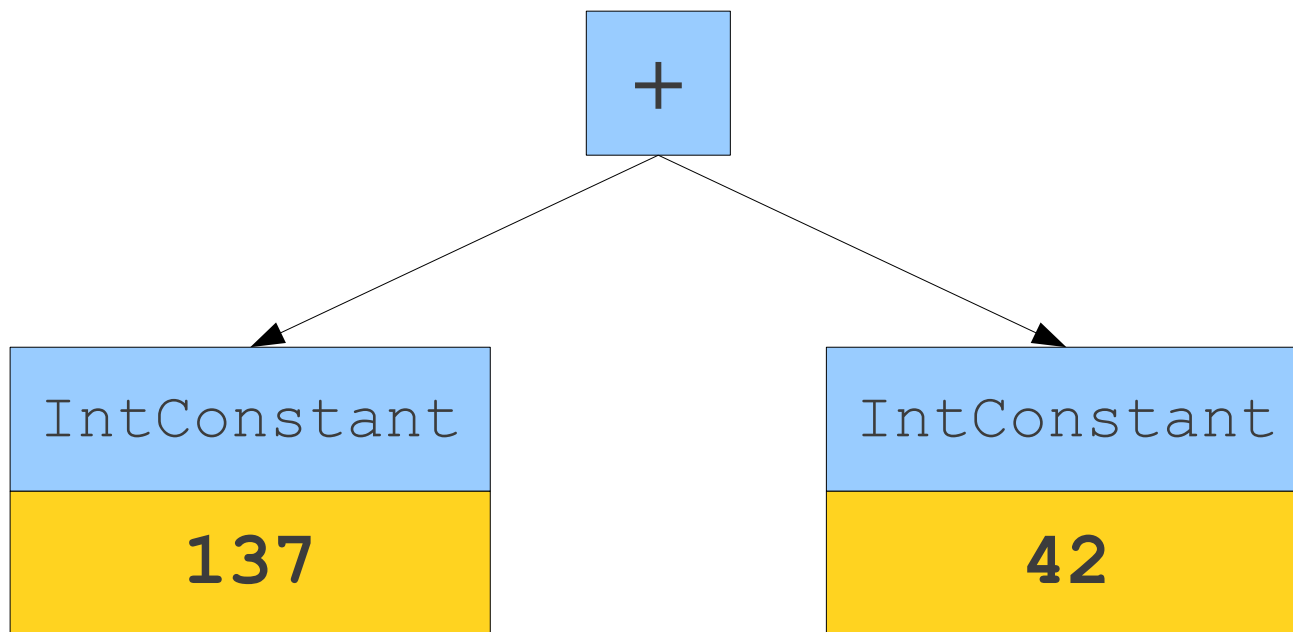
# Inferring Expression Types

- How do we determine the type of an expression?
- Think of process as **logical inference**.



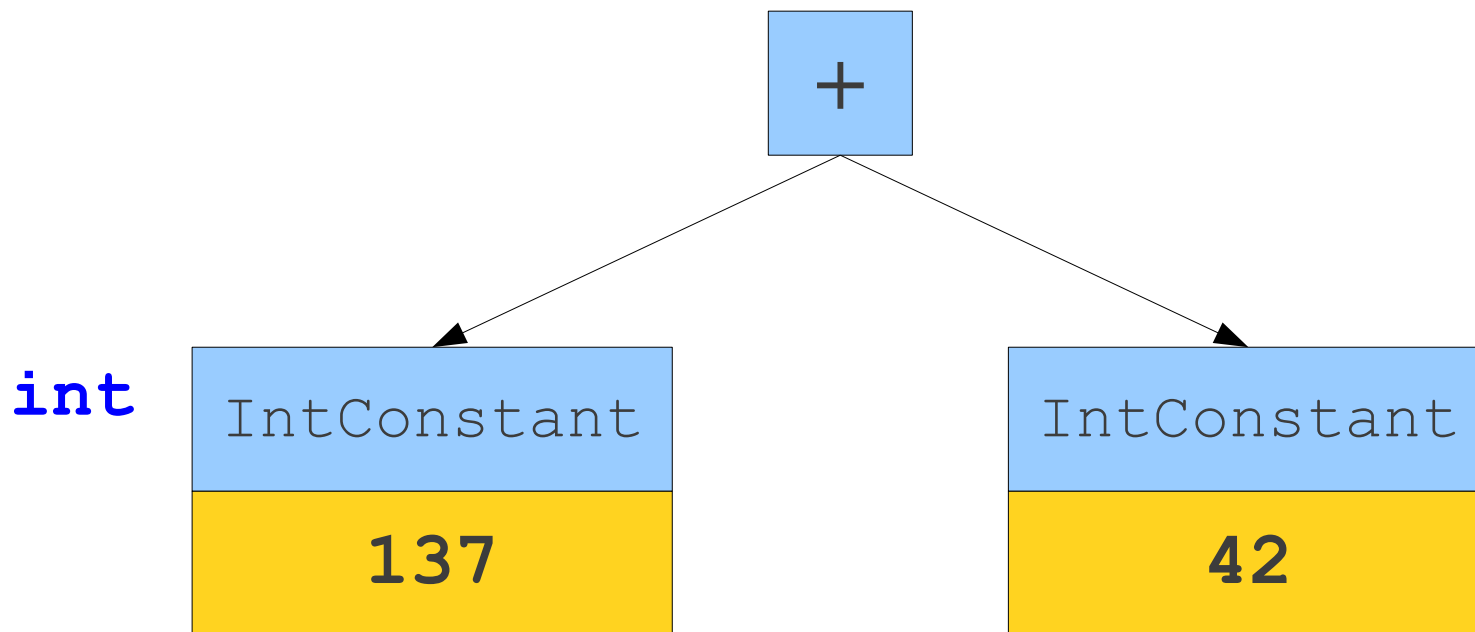
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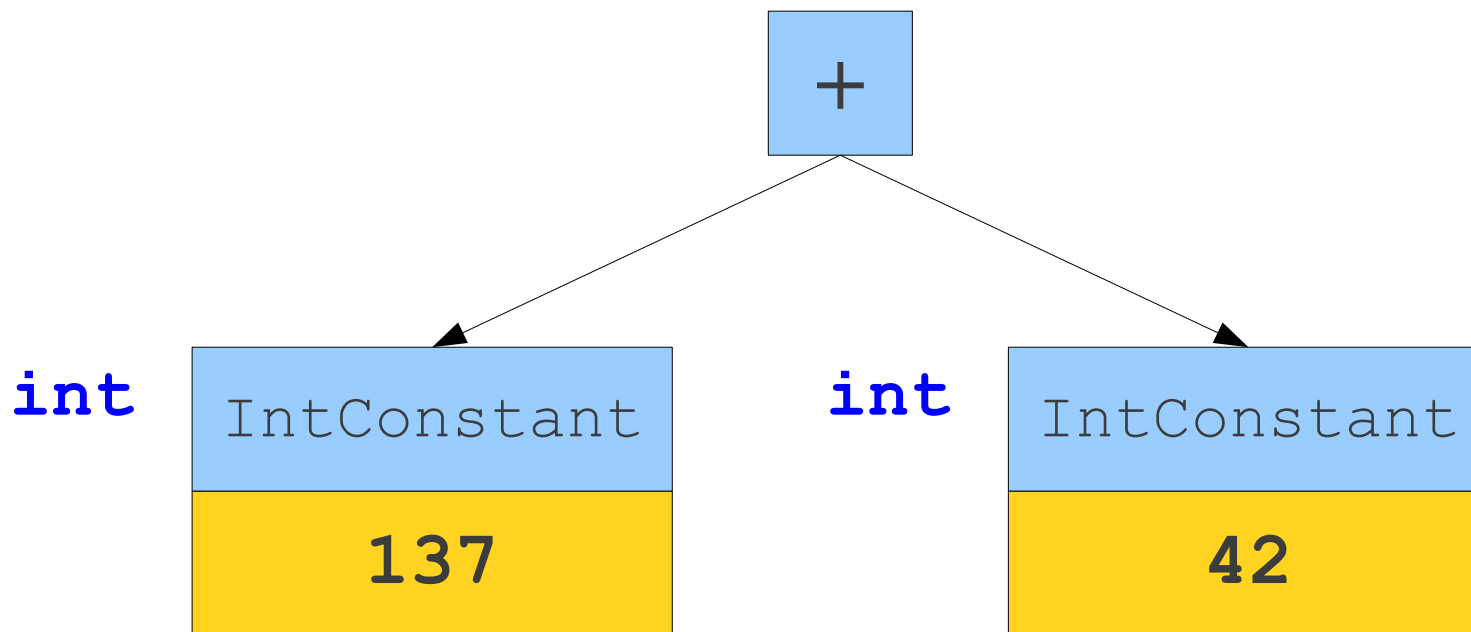
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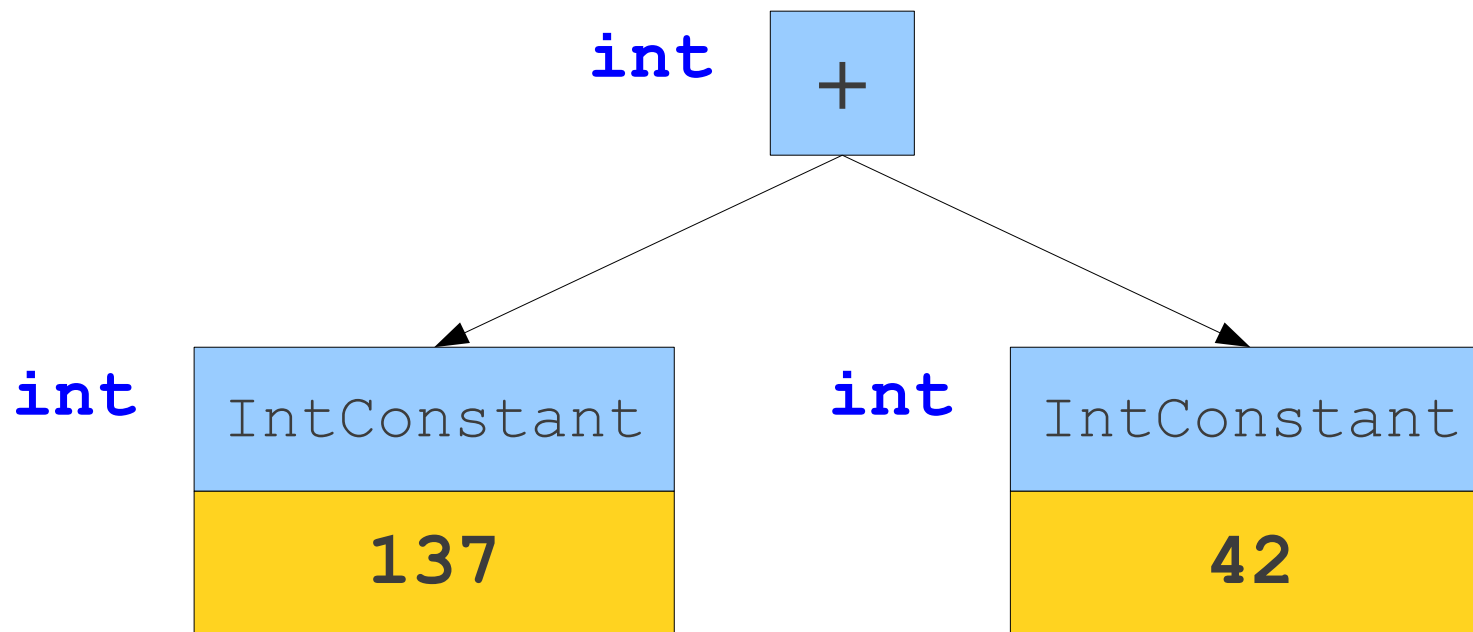
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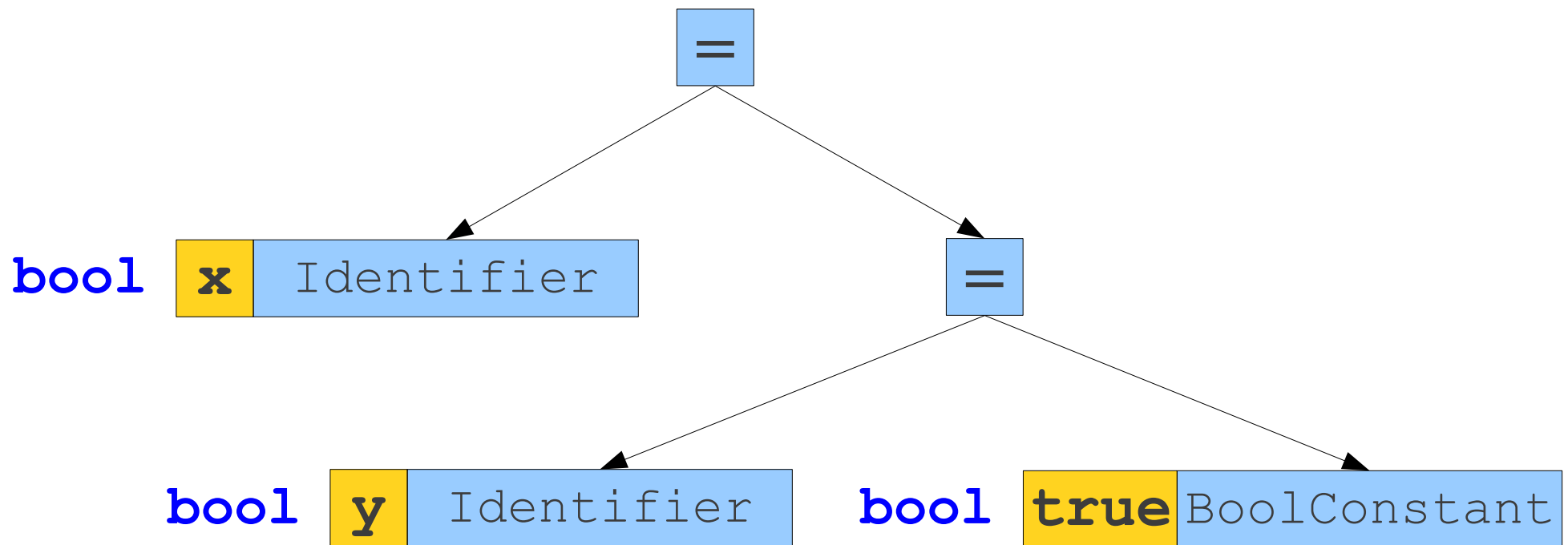


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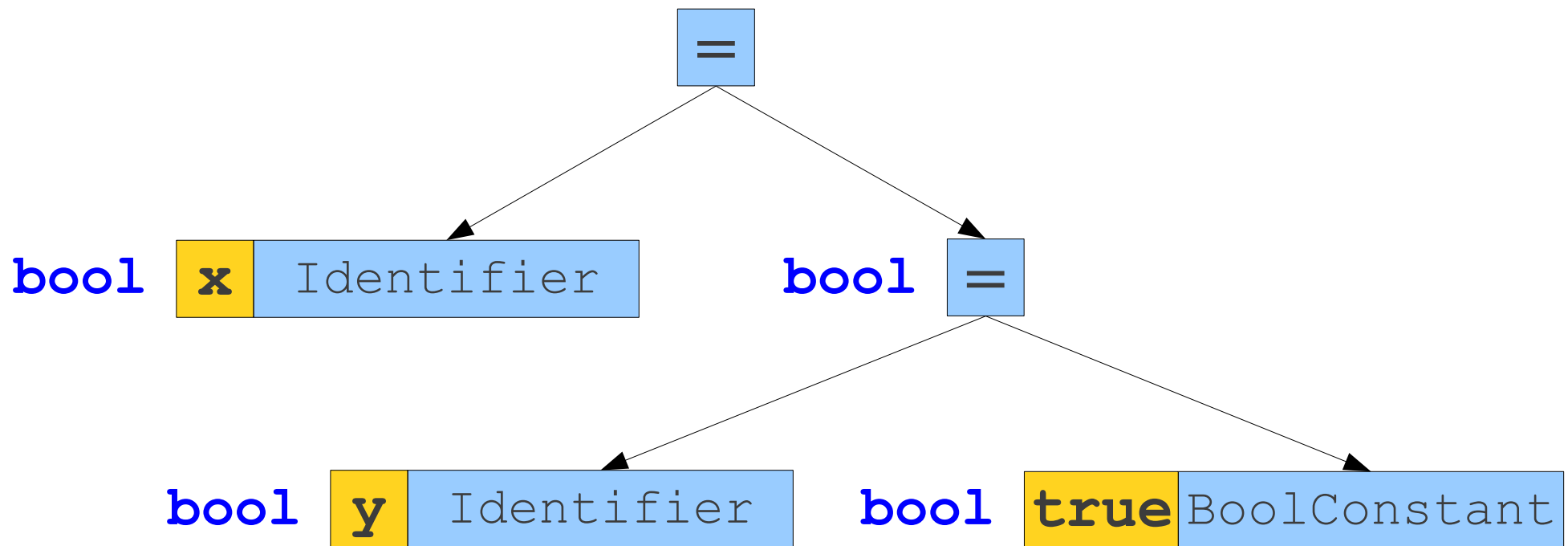
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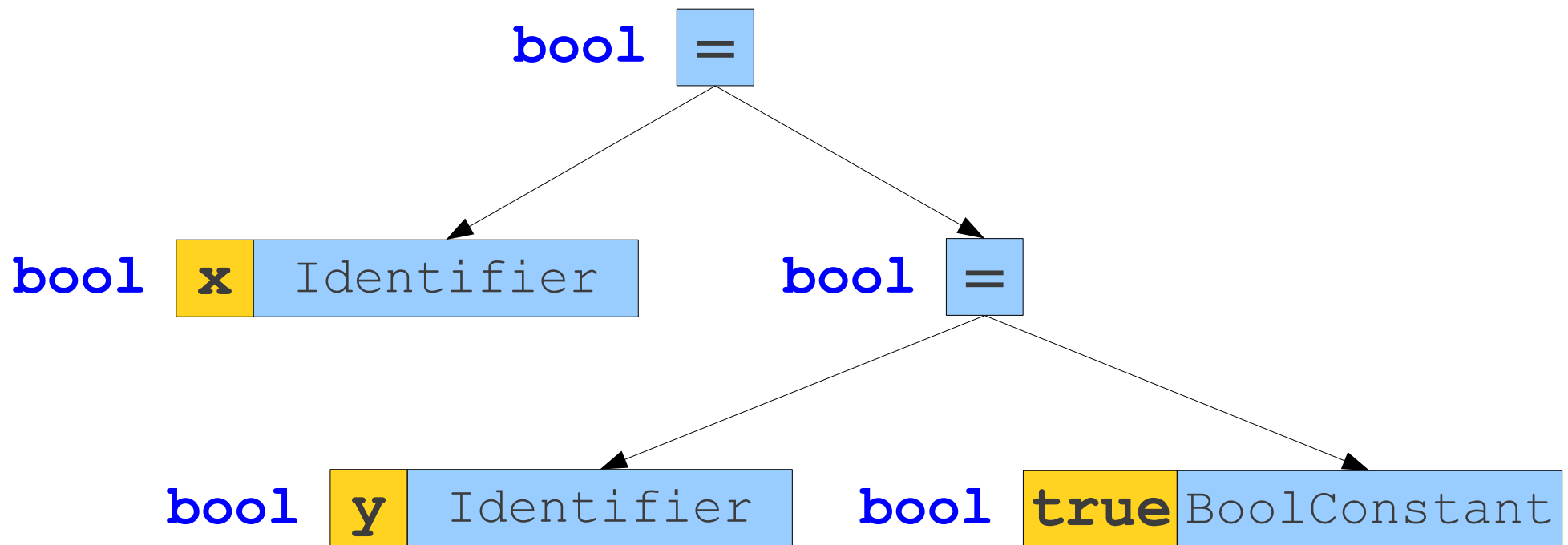
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# Type Checking as Proofs

- We can think of syntax analysis as proving claims about the types of expressions.
- We begin with a set of **axioms**, then apply our **inference rules** to determine the types of expressions.
- Many type systems can be thought of as proof systems.

# Sample Inference Rules

- “If  $x$  is an identifier that refers to an object of type  $t$ , the expression  $x$  has type  $t$ .”
- “If  $e$  is an integer constant,  $e$  has type  $\mathbf{int}$ .”
- “If the operands  $e_1$  and  $e_2$  of  $e_1 + e_2$  are known to have types  $\mathbf{int}$  and  $\mathbf{int}$ , then  $e_1 + e_2$  has type  $\mathbf{int}$ .”